



Also available at http://amc-journal.eu ISSN 1855-3966 (printed edn.), ISSN 1855-3974 (electronic edn.) ARS MATHEMATICA CONTEMPORANEA 9 (2015) 45–50

# Comparing the irregularity and the total irregularity of graphs

Darko Dimitrov

Institut für Informatik, Freie Universität Berlin, 14195 Berlin, Germany

Riste Škrekovski \*

Department of Mathematics, University of Ljubljana, 1000 Ljubljana and Faculty of Information Studies, 8000 Novo Mesto, Slovenia

Received 9 June 2012, accepted 15 August 2013, published online 13 June 2014

#### Abstract

Albertson [4] has defined the *irregularity* of a simple undirected graph G as  $\operatorname{irr}(G) = \sum_{uv \in E(G)} |d_G(u) - d_G(v)|$ , where  $d_G(u)$  denotes the degree of a vertex  $u \in V(G)$ . Recently, in [1] a new measure of irregularity of a graph, so-called the *total irregularity*, was defined as  $\operatorname{irr}_t(G) = \frac{1}{2} \sum_{u,v \in V(G)} |d_G(u) - d_G(v)|$ . Here, we compare the irregularity and the total irregularity of graphs. For a connected graph G with n vertices, we show that  $\operatorname{irr}_t(G) \leq n^2 \operatorname{irr}(G)/4$ . Moreover, if G is a tree, then  $\operatorname{irr}_t(G) \leq (n-2)\operatorname{irr}(G)$ .

*Keywords: The irregularity of graph, the total irregularity of graph. Math. Subj. Class.: 05C05, 05C07,05C99* 

## 1 Introduction

Let G be a simple undirected graph of order n = |V(G)| and size m = |E(G)|. For  $v \in V(G)$ , the degree of v, denoted by  $d_G(v)$ , is the number of edges incident to v. Albertson [4] defines the *imbalance* of an edge  $e = uv \in E(G)$  as  $imb_G(uv) = |d_G(u) - d_G(v)|$ and the *irregularity* of G as

$$\operatorname{irr}(G) = \sum_{uv \in E(G)} \operatorname{imb}_G(uv).$$
(1.1)

Obviously, a connected graph G has irregularity zero if and only if G is regular. In [4] Albertson presented upper bounds on irregularity for bipartite graphs, triangle-free graphs

<sup>\*</sup>Partially supported by Slovenian ARRS Program P1-00383 and Creative Core - FISNM - 3330-13-500033. *E-mail addresses:* dimdar@zedat.fu-berlin.de (Darko Dimitrov), skrekovski@gmail.com (Riste Škrekovski)

and arbitrary graphs, as well as a sharp upper bound for trees. Some results about the irregularity of bipartite graphs are given in [4, 14]. Related to the work of Albertson is the work of Hansen and Mélot [13], who characterized the graphs with n vertices and m edges with maximal irregularity. Various upper bounds on the irregularity of a graph were given in [19], where  $K_{r+1}$ -free graphs, trees and unicyclic graphs with fixed number of vertices of degree one were considered. In [16], relations between the irregularity and the matching number of trees and unicyclic graphs were investigated. More results on irregularity, imbalance and related measures, one can find in [3, 5, 6, 17, 18].

Recently, in [1] a new measure of irregularity of a simple undirected graph, so-called the *total irregularity*, was defined as

$$\operatorname{irr}_{t}(G) = \frac{1}{2} \sum_{u,v \in V(G)} |d_{G}(u) - d_{G}(v)|.$$
(1.2)

Other approaches, that characterize how irregular a graph is, have been proposed [2, 3, 7, 8, 9, 10, 15]. In this paper, we focus on the relation between the irregularity (1.1) and the total irregularity (1.2) of a graph.

In the sequel we introduce the notation used in the rest of the paper. For  $u, v \in V(G)$ , we denote by  $d_G(u, v)$  the length of a shortest path in G between u and v. In this short paper the notation of the sets, that will be defined next, is always regarding the graph G we consider. By  $V_{a,b}$ , we denote a set of vertices of a graph with degrees in [a, b], and by  $V_{\geq a}$  (resp.  $V_{\leq a}$ ), we denote a set of vertices of a graph with degrees at least a (resp. with degrees at most a). Similarly, by  $V_{\geq a}^x$  (resp.  $V_{\leq a}^x$ ), we denote a set of resp.  $v_{\leq a}$ ), we denote a set of neighboring vertices of a vertex x with degrees at least a (resp.  $v_{\leq a}^x$ ), we denote a set of the above mentioned sets, we denote by small v (e.g.,  $v_{\leq a} = |V_{\leq a}|$  or  $v_{\leq a}^x = |V_{\leq a}|$ ).

A subgraph  $T = v_1 v_2 \cdots v_l$  of a graph G, where  $v_l$  is a leaf in G, is called a *tread* if  $d_G(v_1) = d_G(v_2) = \cdots = d_G(v_{l-1}) = 2$ , and  $v_1$  is adjacent to a vertex with degree at least three. Let  $T_1 = v_1 v_2 \cdots v_s$  and  $T_2 = u_1 u_2 \cdots u_l$  be two threads of a graph G with leaves  $v_s$  and  $u_l$ , respectively, and let  $v_0$  be the other neighbour of  $v_1$ . By  $G' = G(T_2 \circ T_1)$  we denote a graph that is obtained from G after a *concatenation* of  $T_2$  to  $T_1$ , i.e., after deleting the edge  $v_0 v_1$  and adding an edge between  $u_l$  and  $v_1$ .

#### 2 General graphs

Obviously,  $\operatorname{irr}(G) \leq \operatorname{irr}_t(G)$ . And, it is not hard to show that equality holds precisely when all non-adjacent vertices have same degree. Such a class of graphs are the complete *k*-partite graphs. More examples of graphs with equal irregularity and total irregularity can be found in [11]. Now, we give an upper bound on  $\operatorname{irr}_t(G)$  in term of  $\operatorname{irr}(G)$ .

**Theorem 2.1.** Let G be a connected graph on n-vertices. Then

$$\operatorname{irr}_t(G) \le \frac{n^2}{4}\operatorname{irr}(G).$$

Moreover, the bound is sharp for infinitely many graphs.

*Proof.* Let T be a spanning tree of G. Then, any two vertices a, b of G are connected by an unique path  $P_{ab} = x_1 x_2 \cdots x_s$  in T, where  $x_1 = a$  and  $x_s = b$ . By the triangle inequality,

we have that

$$\operatorname{irr}_{t}(G) = \frac{1}{2} \sum_{a,b \in V(G)} |d_{G}(a) - d_{G}(b)|$$

$$\leq \frac{1}{2} \sum_{a,b \in V(G)} |d_{G}(x_{1}) - d_{G}(x_{2})| + |d_{G}(x_{2}) - d_{G}(x_{3})| + \cdots$$
(2.1)

For an edge  $uv \in E(T)$ , let  $n_u = \{x \mid x \in V(T) \text{ and } d_T(x, u) < d_T(x, v)\}$ . Similarly, let  $n_v = \{x \mid x \in V(T) \text{ and } d_T(x, u) > d_T(x, v)\}$ . Each summand  $|d_G(u) - d_G(v)|$  in the last sum of (2.1) occurs in the sum exactly  $n_{uv} = n_u n_v$  times. Also, each summand  $|d_G(v) - d_G(u)|$  occurs  $n_{uv}$  times. Thus,

$$\operatorname{irr}_t(G) \leq \sum_{uv \in E(T)} |d_G(u) - d_G(v)| n_{uv}.$$

As  $n_{uv} \leq (n/2)(n/2) = n^2/4$ , and  $\sum_{uv \in E(T)} |d_G(u) - d_G(v)| \leq \sum_{uv \in E(G)} |d_G(u) - d_G(v)|$ , we obtain the desired inequality.

Now, we show that the bound  $n^2/4$  is sharp. Let a, b be two distinct integers, say a < b. Consider a graph  $G_a$  whose all vertices are of degree a, with exception of one vertex u which is of degree a - 1. Similarly, consider a graph  $G_b$  whose all vertices are of degree b, with exception of one vertex u which is of degree b - 1. Let  $G^*$  be the graph obtained from  $G_a$  and  $G_b$  by connecting u and v. Let  $n_a = |V(G_a)|$  and  $n_b = |V(G_b)|$ . Observe that  $\operatorname{irr}(G^*) = b - a$  and  $\operatorname{irr}_t(G^*) = (b - a)n_an_b$ . Choosing  $n_a = n_b = n/2$ , we obtain

$$\frac{\operatorname{irr}_t(G^*)}{\operatorname{irr}(G^*)} = n_a n_b = \frac{n^2}{4}.$$

In order to show that such graphs  $G_a$  and  $G_b$  exist, one may use the theorem of Erdős-Gallai [12] which states that a sequence  $d_1 \ge d_2 \ge \cdots \ge d_n$  of non-negative integers with even sum is graphic (i.e., there exist a graph with such a degree sequence) if and only if

$$\sum_{i=1}^{r} d_i \le r(r-1) + \sum_{i=r+1}^{n} \min(r, d_i),$$
(2.2)

for all  $1 \leq r \leq n$ .

So, fix a, b, and  $n_a = n_b$  to be odd numbers with  $n_a \gg \max\{a, b\}$ . We will show the existence of the graph  $G_a$ . In a similar way, one can show the existence of the graph  $G_b$ . As  $(n_a - 1)a + (a - 1)$  is even, the parity condition of the theorem of Erdős-Gallai is satisfied. So, we need to show only (2.2). For this we consider three cases regarding r and a:

- r ≤ a − 1. Then, (2.2) can be written as ra ≤ r(r − 1) + (n<sub>a</sub> − r)r. It obviously holds since a ≪ n<sub>a</sub> − r.
- r = a. In this case, (2.2) can be written as  $ra \le r(r-1) + (n_a r)r 1$ , which holds for a similar reason as the previous case.
- $r \ge a + 1$ . Similarly, (2.2) can be written as  $ra \le r(r-1) + (n_a r)a 1$ , and it holds as  $ra \ll r(r-1)$ .

### **3** Trees

In this section, we give an upper bound on  $\operatorname{irr}_t(G)$  in term of  $\operatorname{irr}(G)$ , when G is a tree. To show the bound, we will use the following lemma.

**Lemma 3.1.** Let G be a tree, x a vertex of degree  $d \ge 3$  incident with threads  $T_1$  and  $T_2$ , and let  $G' = G(T_2 \circ T_1)$ . Then,

(a) 
$$\operatorname{irr}_t(G) - \operatorname{irr}_t(G') = 2v_{2,d-1};$$
  
(b)  $\operatorname{irr}(G) - \operatorname{irr}(G') = 2(d - v_{\geq d}^x - 1).$ 

*Proof.* Let  $T_1 = a_1 a_2 \cdots a_s$  and  $T_2 = b_1 b_2 \cdots b_l$ . We consider the identities separately.

(a) Notice that all other vertices except x and  $b_l$  have the same degree in G and G'. Hence, it holds that

$$\operatorname{irr}_{t}(G) - \operatorname{irr}_{t}(G') = \sum_{u \neq b_{l}} (|d_{G}(x) - d_{G}(u)| - |d_{G'}(x) - d_{G'}(u)|) + \sum_{u \neq x} (|d_{G}(u) - d_{G}(b_{l})| - |d_{G'}(u) - d_{G'}(b_{l})|) + |d_{G}(x) - d_{G}(b_{l})| - |d_{G'}(x) - d_{G'}(b_{l})|.$$

Since  $d_{G'}(x) = d_G(x) - 1 = d - 1$  and  $d_{G'}(b_l) = d_G(b_l) + 1 = 2$ , further we have

$$\operatorname{irr}_{t}(G) - \operatorname{irr}_{t}(G') = \sum_{u \neq b_{l}} (|d - d_{G}(u)| - |d - 1 - d_{G}(u)|) + \sum_{u \neq x} (|d_{G}(u) - 1| - |d_{G}(u) - 2|) + 2.$$
(3.1)

If  $u \in V_{\leq d-1}$ , then  $|d - d_G(u)| - |d - 1 - d_G(u)| = 1$ , otherwise  $|d - d_G(u)| - |d - 1 - d_G(u)| = -1$ . Hence, the first sum in (3.1) is equal to  $v_{\leq d-1} - 1 - v_{\geq d}$ . Similarly, if  $u \in V_{\geq 2}$ , then  $|d_G(u) - 1| - |d_G(u) - 2| = 1$ , otherwise  $|d_G(u) - 1| - |d_G(u) - 2| = -1$ . Thus, the second sum in (3.1) is equal to  $v_{\geq 2} - 1 - v_1$ . Applying these observations, we have

$$irr_t(G) - irr_t(G') = v_{\leq d-1} - 1 - v_{\geq d} + v_{\geq 2} - 1 - v_1 + 2$$
  
=  $v_{\leq d-1} - v_1 + v_{\geq 2} - v_{\geq d}$   
=  $2v_{2,d-1}$ .

(b) Let e<sub>1</sub> = xa<sub>1</sub>, e<sub>2</sub> = xb<sub>1</sub>, e<sub>3</sub> = b<sub>l-1</sub>b<sub>l</sub> and E<sub>1</sub> = {e<sub>1</sub>, e<sub>2</sub>, e<sub>3</sub>}. Denote by E<sub>2</sub> the set of edges incident to x that are different from e<sub>1</sub> and e<sub>2</sub>. Notice that every edge not in E<sub>1</sub> ∪ E<sub>2</sub> contributes zero to the difference irr(G) - irr(G'). So, we can infer

$$\begin{split} \operatorname{irr}(G) - \operatorname{irr}(G') &= \sum_{uv \in E_2} (\operatorname{imb}_G(uv) - \operatorname{imb}_{G'}(uv)) \\ &+ \sum_{uv \in E_1} (\operatorname{imb}_G(uv) - \operatorname{imb}_{G'}(uv)). \end{split}$$

Notice that the first sum is equal to  $-v_{\geq d}^* + (v_{\leq d-1}^* - 2)$  (we have -2 as the edges  $e_1$ and  $e_2$  are excluded in this sum). In  $\overline{G'}$ , the edge  $e_1 = xa_1$  does not exist anymore, but there is a new edge  $e'_1 = b_l a_1$ . Observe that after the concatenation  $T_2 \circ T_1$  all other edges preserve their end-vertices. First, we consider the contribution of  $e_1$  and  $e'_1$  in  $\operatorname{irr}(G) - \operatorname{irr}(G')$ . There are two possibilities regarding the length of  $T_1$ :

- s = 1: Then,  $\operatorname{imb}_{G}(e_1) = d 1$  and  $\operatorname{imb}_{G'}(e'_1) = 1$ ;
- $s \ge 2$ : In this case,  $\operatorname{imb}_G(e_1) = d 2$  and  $\operatorname{imb}_{G'}(e'_1) = 0$ .

In both of them, we obtain  $\operatorname{imb}_G(e_1) - \operatorname{imb}_{G'}(e'_1) = d - 2$ .

Next, we consider the contributions of  $e_2$  and  $e_3$  together. Again, consider two possibilities regarding the length of  $T_2$ :

- l = 1: Then,  $e_2 = e_3$  and  $\operatorname{imb}_G(e_2) = d 1$  and  $\operatorname{imb}_{G'}(e_2) = d 3$ ;
- $l \ge 2$ : In this case,  $e_2 \ne e_3$ , and  $\operatorname{imb}_G(e_2) = d 2$ ,  $\operatorname{imb}_{G'}(e_2) = d 3$ ,  $\operatorname{imb}_G(e_3) = 1$  and  $\operatorname{imb}_{G'}(e_3) = 0$ .

In both cases, we obtain that  $\sum_{e \in \{e_2, e_3\}} (imb_G(e) - imb_{G'}(e)) = 2$ . So finally, we have that

$$\begin{split} \operatorname{irr}(G) - \operatorname{irr}(G') &= -v_{\geq d}^x + (v_{\leq d-1}^x - 2) + d - 2 + 2 \\ &= -v_{\geq d}^x + v_{\leq d-1}^x - 2 + d \\ &= 2(d - v_{\geq d}^x - 1). \end{split}$$

**Theorem 3.1.** Let G be a tree with n vertices. Then

$$\operatorname{irr}_t(G) \le (n-2)\operatorname{irr}(G).$$

Moroever, equality holds if and only if G is a path.

*Proof.* Let  $n_1(G)$  be the number of vertices of G with degree one. We will prove the second inequality by induction on  $n_1(G)$ . If  $n_1(G) = 0$ , then  $G \simeq P_1$ ,  $irr(G) = irr_t(G) = 0$ , and the equality in the theorem holds. Since G is a tree,  $n_1(G) \neq 1$ . If  $n_1(G) = 2$ , then  $G \simeq P_n$ . In this case irr(G) = 2 and  $irr_t(G) = 2(n-2)$ , hence we obtain equality.

Now, assume  $n_1(G) > 2$ . Then, it is easy to see that G has a vertex x of degree  $d \ge 3$ , incident with at least two threads  $T_1$  and  $T_2$ . Let  $G' = G(T_2 \circ T_1)$ . Since  $n_1(G') = n_1(G) - 1$ , we can assume that inequality holds for G', i.e.,

$$\operatorname{irr}_t(G') \le (n-2)\operatorname{irr}(G'). \tag{3.2}$$

By Lemma 3.1, we have

$$\operatorname{irr}(G') = \operatorname{irr}(G) - 2(d - v_{\geq d}^x - 1)$$
 and  $\operatorname{irr}_t(G') = \operatorname{irr}_t(G) - 2v_{2,d-1}$ . (3.3)

Plugging (3.3) in (3.2), we obtain

$$(n-2)\operatorname{irr}(G) \ge \operatorname{irr}_t(G) - 2v_{2,d-1} + 2(n-2)(d-v_{\ge d}^x - 1).$$
 (3.4)

As  $d(x) = d \ge 3$  and x is incident with two threads, we infer  $v_{\ge d}^x + 2 \le d$ , and so  $2(d - v_{\ge d}^x - 1) \ge 2$ . Observe also that  $v_{2,d-1} \le n-3$ . Hence  $2(n-2)(d - v_{\ge d}^x - 1) > 2(n-3) \ge 2v_{2,d-1}$ . This together with (3.4) gives  $(n-2)\operatorname{irr}_{d}(G) > \operatorname{irr}_{t}(G)$ .  $\Box$ 

#### References

- H. Abdo, S. Brandt and D. Dimitrov, The total irregularity of a graph, *Discrete Math. Theor. Comput. Sci.* 16 (2014), 201–206.
- [2] Y. Alavi, A. Boals, G. Chartrand, P. Erdős and O. R. Oellermann, k-path irregular graphs, Congr. Numer. 65 (1988) 201–210.
- [3] Y. Alavi, G. Chartrand, F. R. K. Chung, P. Erdős, R. L. Graham and O. R. Oellermann, Highly irregular graphs, J. Graph Theory 11 (1987) 235–249.
- [4] M. O. Albertson, The irregularity of a graph, Ars Comb. 46 (1997) 219–225.
- [5] F. K. Bell, A note on the irregularity of graphs, Linear Algebra Appl. 161 (1992) 45-54.
- [6] Y. Caro and R. Yuster, Graphs with large variance, Ars Comb. 57 (2000) 151-162.
- [7] G. Chartrand, P. Erdős and O. R. Oellermann, How to define an irregular graph, *Coll. Math. J.* 19 (1988) 36–42.
- [8] G. Chartrand, K. S. Holbert, O. R. Oellermann and H. C. Swart, F-degrees in graphs, Ars Comb. 24 (1987) 133–148.
- [9] L. Collatz and U. Sinogowitz, Spektren endlicher Graphen, *Abh. Math. Sem. Univ. Hamburg* 21 (1957) 63–77.
- [10] D. Cvetković and P. Rowlinson, On connected graphs with maximal index, *Publications de l'Institut Mathematique (Beograd)* 44 (1988) 29–34.
- [11] D. Dimitrov, T. Réti, Graphs with equal irregularity indices, *Acta Polytech. Hung.* 11 (2014), 41–57.
- [12] P. Erdős and T. Gallai, Graphs with prescribed degrees of vertices, (in Hungarian) *Mat. Lapok.* 11 (1960) 264–274.
- [13] P. Hansen and H. Mélot, Variable neighborhood search for extremal graphs 9. Bounding the irregularity of a graph, DIMACS Ser. Discrete Math. Theoret. Comput. Sci. 69 (1962) 253–264.
- [14] M. A. Henning and D. Rautenbach, On the irregularity of bipartite graphs, *Discrete Math.* 307 (2007) 1467–1472.
- [15] D. E. Jackson and R. Entringer, Totally segregated graphs, *Congress. Numer.* 55 (1986) 159– 165.
- [16] W. Luo and B. Zhou, On the irregularity of trees and unicyclic graphs with given matching number, Util. Math. 83 (2010) 141–147.
- [17] D. Rautenbach and I. Schiermeyer, Extremal problems for imbalanced edges, *Graphs Comb.* 22 (2006) 103–111.
- [18] D. Rautenbach and L. Volkmann, How local irregularity gets global in a graph, J. Graph Theory 41 (2002) 18–23.
- [19] B. Zhou and W. Luo, On irregularity of graphs, Ars Combin. 88 (2008) 55-64.