

# The parameters of Fibonacci and Lucas cubes\*

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## Abstract

Motivated by the conjectures from Castro, et al. in 2011, in this paper we use integer programming formulations for computing the domination number, the 2-packing number and the independent domination number of Fibonacci cubes and Lucas cubes for  $n \leq 13$ .

*Keywords:* Fibonacci cubes, Lucas cubes, domination number, 2-packing number.

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## 1 Introduction

Hypercubes form one of the most applicable classes of graphs with many appealing properties. The  $n$ -cube  $Q_n$  is the graph whose vertices are all binary strings of length  $n$ , and two vertices are adjacent if they differ in exactly one position. The Fibonacci cubes were introduced as a model for interconnection networks [4, 2]. They offer challenging mathematical and computational problems, and admit a recursive decomposition into smaller Fibonacci cubes (see [5], [6], [8] for their structural properties). The Fibonacci cubes can be recognized in  $O(m \log n)$  time (where  $n$  is the order and  $m$  the size of a given graph) [10]. The Lucas cubes [7] form a class of graphs closely related to the Fibonacci cubes, obtained by removing some vertices from the Fibonacci cubes.

Let  $Q_n$  be the  $n$ -dimensional hypercube. A Fibonacci string of length  $n$  is a binary string  $b_1 b_2 \dots b_n$  with  $b_i \cdot b_{i+1} = 0$  for  $1 \leq i < n$ . In other words, Fibonacci strings are binary strings that contain no consecutive ones. The Fibonacci cube  $\Gamma_n$ , for  $n \geq 1$  is the subgraph of  $Q_n$  induced by the Fibonacci strings of length  $n$ . A Fibonacci string  $b_1 b_2 \dots b_n$  is a Lucas string if  $b_1 \cdot b_n = 0$ . In other words, Lucas strings are binary strings

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that contain no consecutive ones circularly. The Lucas cube  $\Lambda_n$ , for  $n \geq 1$  is the subgraph of  $Q_n$  induced by the Lucas strings of length  $n$ . It is well-known that  $|V(\Gamma_n)| = F_{n+2}$ , where  $F_n$  are the Fibonacci numbers:  $F_0 = 0, F_1 = 1, F_{n+1} = F_n + F_{n-1}$  for  $n \geq 1$ . Similarly,  $|V(\Lambda_n)| = L_n$  for  $n \geq 1$ , where  $L_n$  are the Lucas numbers:  $L_0 = 2, L_1 = 1, L_{n+1} = L_n + L_{n-1}$  for  $n \geq 1$ .

Let  $G$  be a graph. Set  $D \subseteq V(G)$  is a dominating set if every vertex from  $V(G)$  either belongs to  $D$  or is adjacent to some vertex from  $D$ . The domination number  $\gamma(G)$  is the minimum cardinality of a dominating set of  $G$ . A set  $X \subseteq V(G)$  is called a 2-packing if  $d(u, v) > 2$  for any two different vertices  $u$  and  $v$  of  $X$ . The 2-packing number  $\rho(G)$  is the maximum cardinality of a 2-packing of  $G$ . It is well-known that for any graph  $G$  holds  $\gamma(G) \geq \rho(G)$ .

An independent set or stable set is a set of vertices in a graph, no two of which are adjacent. The independent domination number  $i(G)$  of a graph  $G$  is the size of the smallest independent dominating set (or, equivalently, the size of the smallest maximal independent set). The minimum dominating set in a graph will not necessarily be independent, but the size of a minimum dominating set is always less than or equal to the size of a minimum maximal independent set,  $\gamma(G) \leq i(G)$ .

Pike and Zou in [9] obtained a lower bound for the domination number of Fibonacci cube of order  $n$  and determined the exact value of the domination number of Fibonacci cubes of order at most 8. Castro et al. in [1] obtained upper and lower bounds for the domination and 2-packing number of Fibonacci and Lucas cubes. Furthermore, the authors obtained the exact values for  $\gamma(\Gamma_n)$  and  $\gamma(\Lambda_n)$  for  $n \leq 9$  and for  $\rho(\Gamma_n)$  and  $\rho(\Lambda_n)$  for  $n \leq 10$ .

In this paper we use integer programming method to compute the exact values of the domination, 2-packing and independent domination number of Fibonacci and Lucas cubes for  $n \leq 13$ , which resolves the conjecture from [1].

## 2 Main results

For each subset of the vertex set  $S \subseteq V(G)$  define

$$x_i = \begin{cases} 1 & \text{if } i \in S \\ 0 & \text{if } i \in V \setminus S. \end{cases}$$

The neighborhood  $N(v)$  of a vertex  $v$  in a graph  $G$  is the induced subgraph of  $G$  consisting of all vertices adjacent to  $v$  and all edges connecting two such vertices. Let  $N[v] = N(v) \cup \{v\}$  denote the closed neighborhood of the vertex  $v$ .

The domination number of  $G$  can be formulated as the following 0–1 integer programming problem:

$$\gamma(G) = \min \sum_{i=1}^n x_i \tag{2.1}$$

subject to

$$\sum_{j \in N[i]} x_j \geq 1, \tag{2.2}$$

$$x_i \in \{0, 1\}, \quad \text{for all } 1 \leq i \leq n. \tag{2.3}$$

It is easy to see that the conditions (2.2) and (2.3) define dominating set  $S$  and vice versa [3]. For Fibonacci cube  $\Gamma_n$  this formulation has  $F_{n+2}$  variables and  $2F_{n+2}$  constraints,

while each condition from (2.2) contains at most  $n$  variables. For Lucas cube  $\Lambda_n$  this formulation has  $L_n$  variables and  $2L_n$  constraints, while each condition from (2.2) contains at most  $n$  variables.

The 2-packing number of  $G$  can be formulated as the following 0 – 1 integer programming problem:

$$\rho(G) = \max \sum_{i=1}^n x_i \quad (2.4)$$

subject to

$$\sum_{j \in N[i]} x_j \leq 1, \quad (2.5)$$

$$x_i \in \{0, 1\}, \quad \text{for all } 1 \leq i \leq n. \quad (2.6)$$

We will prove that the conditions (2.5) and (2.6) define 2-packing set  $S$  and vice versa. Let  $S$  be a 2-packing set. Since  $S$  does not contain two vertices on distance 1 or 2, for each  $v \in V(G)$  there is at most one vertex from the closed neighborhood  $N[v]$  which belongs to  $S$ . Assume now that the set  $S$  satisfies the condition (2.5) and let  $u$  and  $v$  be two vertices from  $S$  on distance 2. In that case for the shortest path  $vwu$ , we have  $\sum_{j \in N[w]} x_j \geq 2$ , which is impossible. Therefore,  $S$  is a 2-packing set.

The independent domination number  $G$  can be formulated as the following 0 – 1 integer programming problem:

$$i(G) = \min \sum_{i=1}^n x_i \quad (2.7)$$

subject to

$$\sum_{j \in N[i]} x_j \geq 1, \quad (2.8)$$

$$(n - 1)x_i + \sum_{j \in N(i)} x_j \leq n - 1, \quad (2.9)$$

$$x_i \in \{0, 1\}, \quad \text{for all } 1 \leq i \leq n. \quad (2.10)$$

The conditions (2.8) and (2.10) define domination set  $S$ , while the condition (2.9) ensures the independence. For  $x_i = 0$  we have always true  $\sum_{j \in N(i)} x_j \leq n - 1$ , while for  $x_i = 1$  we have  $\sum_{j \in N(i)} x_j \leq 0$  which is equivalent to  $\sum_{j \in N[i]} x_j = 1$ . This proves that the formulation is correct. For Fibonacci cube  $\Gamma_n$  this formulation has  $F_{n+2}$  variables and  $3F_{n+2}$  constraints, while each conditions from (2.8) and (2.9) contain at most  $n$  variables. For Lucas cube  $\Lambda_n$  this formulation has  $L_n$  variables and  $3L_n$  constraints, while each condition from (2.8) and (2.9) contain at most  $n$  variables.

The tests were performed on the Intel Core 2 Duo T5800 2.0 GHz with 2 GB RAM running the Linux operating system and using CPLEX 8.1. The results are summarized in Tables 1 and 2. In Tables 3 and 4 we give some examples of dominating sets and 2-packing sets that were obtained during the computation of these values.

These results resolve the conjecture from [1] and support Problem 5.1 for  $n \leq 12$ .

$n$	1	2	3	4	5	6	7	8	9	10	11
$ V(\Gamma_n) $	2	3	5	8	13	21	34	55	89	144	233
$ E(\Gamma_n) $	1	2	5	10	20	38	71	130	235	420	744
$\gamma(\Gamma_n)$	1	1	2	3	4	5	8	12	17	25	
$\rho(\Gamma_n)$	1	1	2	2	3	5	6	9	14	20	29
$i(\Gamma_n)$	1	1	2	3	4	5	8	12	19	26	

Table 1: Parameters of small Fibonacci cubes.

$n$	1	2	3	4	5	6	7	8	9	10	11	12
$ V(\Lambda_n) $	1	3	4	7	11	18	29	47	76	123	199	322
$ E(\Lambda_n) $	0	2	3	8	15	30	56	104	189	340	605	1068
$\gamma(\Lambda_n)$	1	1	1	3	4	5	7	11	16	23	35	
$\rho(\Lambda_n)$	1	1	1	2	3	5	6	8	13	18	26	38
$i(\Lambda_n)$	1	1	1	3	4	5	8	11	17	24	35	

Table 2: Parameters of small Lucas cubes.

Dominating set	
$\Gamma(10)$	$\Lambda(11)$
(0, 1, 0, 1, 0, 0, 0, 0, 0, 0), (0, 1, 0, 0, 0, 1, 0, 0, 0, 0) (1, 0, 1, 0, 0, 1, 0, 0, 0, 0), (1, 0, 0, 1, 0, 0, 1, 0, 0, 0) (0, 0, 0, 0, 1, 0, 1, 0, 0, 0), (0, 0, 0, 0, 0, 0, 0, 1, 0, 0) (1, 0, 1, 0, 1, 0, 0, 1, 0, 0), (1, 0, 0, 1, 0, 1, 0, 1, 0, 0) (1, 0, 0, 0, 0, 0, 0, 0, 1, 0), (0, 0, 1, 0, 0, 0, 0, 0, 1, 0) (0, 1, 0, 0, 1, 0, 0, 0, 1, 0), (0, 0, 0, 1, 0, 1, 0, 0, 1, 0) (0, 0, 1, 0, 0, 0, 1, 0, 1, 0), (0, 1, 0, 1, 0, 0, 1, 0, 1, 0) (1, 0, 1, 0, 1, 0, 1, 0, 1, 0), (0, 0, 0, 1, 0, 0, 0, 0, 0, 1) (1, 0, 0, 0, 1, 0, 0, 0, 0, 1), (0, 0, 1, 0, 1, 0, 0, 0, 0, 1) (1, 0, 0, 0, 0, 1, 0, 0, 0, 1), (0, 1, 0, 0, 0, 0, 1, 0, 0, 1) (1, 0, 1, 0, 0, 0, 1, 0, 0, 1), (1, 0, 0, 0, 0, 0, 0, 1, 0, 1) (0, 1, 0, 0, 1, 0, 0, 1, 0, 1), (0, 0, 1, 0, 0, 0, 1, 0, 1, 0) (0, 1, 0, 1, 0, 1, 0, 1, 0, 1)	(1, 0, 1, 0, 1, 0, 0, 0, 0, 0, 0), (1, 0, 0, 0, 0, 1, 0, 0, 0, 0, 0) (0, 1, 0, 0, 0, 0, 1, 0, 0, 0, 0), (0, 0, 0, 1, 0, 0, 1, 0, 0, 0, 0) (0, 0, 1, 0, 1, 0, 1, 0, 0, 0, 0), (0, 0, 0, 0, 0, 0, 0, 1, 0, 0, 0) (1, 0, 1, 0, 1, 0, 0, 1, 0, 0, 0), (0, 0, 0, 1, 0, 1, 0, 1, 0, 0, 0) (1, 0, 0, 1, 0, 0, 0, 0, 1, 0, 0), (0, 0, 0, 0, 1, 0, 0, 0, 1, 0, 0) (1, 0, 0, 1, 0, 0, 0, 0, 0, 1, 0), (0, 0, 0, 1, 0, 0, 0, 0, 1, 0, 0) (0, 0, 1, 0, 0, 1, 0, 0, 1, 0, 0), (0, 1, 0, 0, 1, 0, 0, 0, 0, 1, 0, 0) (0, 0, 1, 0, 0, 1, 0, 0, 1, 0, 0), (0, 1, 0, 1, 0, 0, 0, 0, 0, 1, 0, 0) (1, 0, 1, 0, 0, 0, 1, 0, 0, 1, 0), (0, 0, 0, 1, 0, 0, 1, 0, 0, 1, 0, 0) (0, 0, 1, 0, 0, 0, 1, 0, 0, 1, 0), (0, 1, 0, 1, 0, 0, 1, 0, 0, 1, 0, 0) (1, 0, 1, 0, 0, 0, 0, 0, 0, 1, 0), (0, 1, 0, 1, 0, 0, 0, 0, 0, 1, 0, 0) (0, 1, 0, 0, 0, 1, 0, 0, 0, 0, 1), (0, 1, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0) (0, 1, 0, 0, 0, 1, 0, 0, 0, 0, 1), (0, 1, 0, 0, 1, 0, 0, 0, 0, 0, 1, 0) (0, 0, 1, 0, 0, 0, 0, 0, 0, 0, 1), (0, 1, 0, 0, 1, 0, 0, 0, 0, 0, 1, 0) (0, 1, 0, 0, 0, 0, 1, 0, 0, 0, 1), (0, 1, 0, 0, 0, 1, 0, 0, 0, 0, 1, 0) (0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1), (0, 1, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0) (0, 1, 0, 0, 0, 0, 1, 0, 0, 0, 1), (0, 1, 0, 0, 0, 1, 0, 0, 0, 0, 1, 0) (0, 0, 1, 0, 0, 0, 0, 0, 0, 0, 1), (0, 1, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0) (0, 1, 0, 0, 0, 1, 0, 0, 0, 0, 1), (0, 1, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0) (0, 1, 0, 0, 0, 1, 0, 0, 0, 0, 1), (0, 1, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0) (0, 1, 0, 0, 0, 1, 0, 0, 0, 0, 1), (0, 1, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0) (0, 0, 1, 0, 0, 0, 0, 0, 0, 0, 1), (0, 0, 0, 1, 0, 0, 0, 0, 0, 0, 1, 0) (0, 1, 0, 0, 0, 1, 0, 0, 0, 0, 1), (0, 0, 0, 1, 0, 0, 0, 0, 0, 0, 1, 0) (0, 0, 0, 0, 0, 0, 1, 0, 0, 0, 1), (0, 0, 0, 1, 0, 0, 0, 0, 0, 0, 1, 0) (0, 1, 0, 0, 0, 1, 0, 0, 0, 0, 1), (0, 0, 0, 1, 0, 0, 0, 0, 0, 0, 1, 0) (0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1), (0, 0, 0, 1, 0, 0, 0, 0, 0, 0, 1, 0)

Table 3: Examples of minimal dominating sets for  $\Gamma(10)$  and  $\Lambda(11)$

2-packaging set	
$\Gamma(11)$	$\Lambda(12)$
(0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0), (1, 0, 1, 0, 0, 0, 0, 1, 0, 0, 0) (0, 1, 0, 0, 0, 1, 0, 1, 0, 0, 0), (0, 1, 0, 0, 1, 0, 0, 0, 0, 1, 0, 0) (1, 0, 1, 0, 1, 0, 0, 0, 1, 0, 0), (0, 0, 0, 1, 0, 1, 0, 0, 1, 0, 0) (0, 0, 1, 0, 0, 0, 0, 1, 0, 1, 0), (0, 1, 0, 1, 0, 0, 1, 0, 0, 1, 0, 0) (1, 0, 0, 0, 1, 0, 0, 0, 0, 1, 0), (0, 0, 1, 0, 0, 1, 0, 0, 0, 1, 0, 0) (0, 1, 0, 1, 0, 1, 0, 0, 0, 1, 0), (0, 1, 0, 0, 0, 0, 1, 0, 0, 0, 1, 0) (1, 0, 0, 1, 0, 0, 1, 0, 0, 1, 0), (0, 0, 1, 0, 1, 0, 1, 0, 0, 0, 1, 0) (0, 0, 0, 1, 0, 0, 0, 1, 0, 1, 0), (0, 1, 0, 0, 1, 0, 0, 1, 0, 0, 1, 0) (1, 0, 0, 0, 0, 1, 0, 1, 0, 1, 0), (0, 0, 1, 0, 1, 0, 0, 0, 0, 0, 1, 0) (1, 0, 1, 0, 0, 0, 1, 0, 0, 0, 1), (0, 0, 0, 1, 0, 0, 1, 0, 0, 0, 0, 1) (0, 1, 0, 0, 1, 0, 1, 0, 0, 0, 1), (0, 1, 0, 1, 0, 0, 0, 1, 0, 0, 0, 1) (1, 0, 0, 0, 1, 0, 0, 1, 0, 0, 1), (0, 0, 1, 0, 0, 0, 1, 0, 1, 0, 0, 1) (1, 0, 0, 0, 1, 0, 0, 1, 0, 0, 1), (0, 1, 0, 0, 0, 1, 0, 1, 0, 0, 0, 1) (1, 0, 0, 0, 1, 0, 1, 0, 0, 0, 1), (1, 0, 0, 1, 0, 0, 0, 0, 0, 1, 0, 1) (0, 1, 0, 0, 0, 1, 0, 0, 1, 0, 1), (1, 0, 1, 0, 0, 0, 1, 0, 0, 0, 1, 0, 1) (1, 0, 0, 0, 1, 0, 1, 0, 1, 0, 1)	(0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0), (1, 0, 0, 1, 0, 1, 0, 0, 0, 0, 0, 0) (0, 1, 0, 0, 1, 0, 1, 0, 0, 0, 0, 0), (1, 0, 1, 0, 1, 0, 0, 1, 0, 0, 0, 0) (0, 0, 1, 0, 0, 1, 0, 1, 0, 0, 0, 0), (1, 0, 0, 0, 1, 0, 0, 0, 1, 0, 0, 0) (0, 1, 0, 1, 0, 1, 0, 0, 1, 0, 0, 0), (1, 0, 1, 0, 0, 0, 1, 0, 1, 0, 0, 0) (0, 0, 0, 1, 0, 0, 1, 0, 1, 0, 0, 0), (1, 0, 1, 0, 0, 0, 0, 0, 0, 1, 0, 0) (0, 1, 0, 0, 0, 1, 0, 0, 0, 0, 1, 0), (1, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0, 0) (0, 1, 0, 0, 0, 1, 0, 0, 0, 0, 1, 0), (1, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0, 0) (0, 0, 1, 0, 0, 1, 0, 0, 1, 0, 0, 0), (0, 1, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0, 0) (0, 1, 0, 0, 0, 1, 0, 0, 0, 1, 0, 0), (1, 0, 0, 0, 0, 1, 0, 0, 1, 0, 0, 0) (0, 0, 0, 1, 0, 0, 1, 0, 1, 0, 0), (0, 1, 0, 0, 0, 0, 1, 0, 1, 0, 0, 0) (0, 0, 0, 0, 1, 0, 0, 0, 1, 0, 1, 0), (1, 0, 0, 0, 0, 1, 0, 1, 0, 1, 0, 0) (0, 0, 0, 0, 1, 0, 0, 0, 1, 0, 1, 0), (0, 1, 0, 0, 0, 1, 0, 0, 0, 0, 1, 0, 0) (0, 1, 0, 1, 0, 0, 0, 0, 0, 0, 1, 0), (1, 0, 1, 0, 0, 1, 0, 0, 0, 0, 0, 1, 0) (0, 0, 0, 0, 1, 0, 0, 0, 0, 0, 1, 0), (1, 0, 0, 0, 0, 1, 0, 1, 0, 0, 0, 1, 0) (0, 0, 0, 0, 1, 0, 0, 0, 0, 0, 1, 0), (1, 0, 0, 0, 0, 1, 0, 0, 0, 0, 0, 1, 0) (0, 1, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0), (1, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0) (0, 1, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0), (1, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0) (0, 0, 1, 0, 0, 0, 0, 0, 0, 0, 1), (0, 0, 0, 1, 0, 0, 0, 0, 0, 0, 1, 0, 0) (0, 1, 0, 0, 0, 1, 0, 0, 0, 0, 1), (0, 1, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0, 0) (0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1), (0, 0, 0, 1, 0, 0, 0, 0, 0, 0, 0, 1, 0) (0, 1, 0, 0, 0, 0, 0, 0, 0, 0, 1), (0, 0, 0, 1, 0, 0, 0, 0, 0, 0, 0, 1, 0) (0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1), (0, 0, 0, 1, 0, 0, 0, 0, 0, 0, 0, 1, 0) (0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1), (0, 0, 0, 1, 0, 0, 0, 0, 0, 0, 0, 1, 0) (0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1), (0, 0, 0, 1, 0, 0, 0, 0, 0, 0, 0, 1, 0) (0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1), (0, 0, 0, 1, 0, 0, 0, 0, 0, 0, 0, 1, 0)

Table 4: Examples of 2-packaging sets for  $\Gamma(11)$  and  $\Lambda(12)$

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