

Lobe, edge, and arc transitivity of graphs of connectivity 1

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Abstract

We give necessary and sufficient conditions for lobe-transitivity of locally finite and locally countable graphs whose connectivity equals 1. We show further that, given any biconnected graph Λ and a “code” assigned to each orbit of $\text{Aut}(\Lambda)$, there exists a unique lobe-transitive graph Γ of connectivity 1 whose lobes are copies of Λ and is consistent with the given code at every vertex of Γ . These results lead to necessary and sufficient conditions for a graph of connectivity 1 to be edge-transitive and to be arc-transitive. Countable graphs of connectivity 1 the action of whose automorphism groups is, respectively, vertex-transitive, primitive, regular, Cayley, and Frobenius had been previously characterized in the literature.

Keywords: Lobe, lobe-transitive, edge-transitive, orbit, connectivity.

Math. Subj. Class.: 05C25, 05C63, 05C38, 20B27

1 Introduction

Throughout this article, Γ denotes a connected simple graph. Consider the equivalence relation \cong on the edge-set $E\Gamma$ of Γ whereby $e_1 \cong e_2$ whenever the edges e_1 and e_2 lie on a common cycle of Γ . A *lobe* is a subgraph of Γ induced by an equivalence class with respect to \cong . Equivalently, a *lobe* is a subgraph that either consists of a cut-edge with its two incident vertices or is a maximal biconnected subgraph¹. A vertex is a *cut-vertex* if it belongs to at least two different lobes. Connected graphs other than K_2 have connectivity 1 if and only if they have a cut-vertex. Clearly no *finite* vertex-transitive graph admits a cut-vertex.

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¹The term “lobe” is due to O. Ore [6]. We eschew the term “block” for this purpose, as it leads to ambiguity when discussing imprimitivity.

Graphs of connectivity 1 whose automorphism groups have certain given properties have been characterized. Those whose automorphism groups are, respectively, vertex-transitive, primitive, and regular were characterized in [5]. In particular, primitive planar graphs of connectivity 1 were characterized in [11]². Cayley graphs of connectivity 1 were characterized in [9]. Graphs of connectivity 1 with Frobenius automorphism groups were characterized in [10]. In the present work, we complete this investigation; we characterize graphs of connectivity 1 whose automorphism groups act transitively on their set of lobes. As a consequence, we obtain characterizations of edge-transitive graphs and arc-transitive graphs of connectivity 1.

The conditions for a graph of connectivity 1 to be lobe-transitive or to be vertex-transitive are independent; such a graph may have either property or neither one or both. Such is not the case for edge- and arc-transitivity. In Section 3 we give necessary and sufficient conditions for a graph to be lobe-transitive. We further show that, given any biconnected graph Λ and a “list” of orbit-multiplicities of copies of $\text{Aut}(\Lambda)$, one can construct a lobe-transitive graph of connectivity 1 all of whose lobes are isomorphic to Λ and locally respects the given list. We give necessary and sufficient conditions for a countable graph of connectivity 1 to be edge-transitive in Section 4 and to be arc-transitive in Section 5. As the sets of conditions for these latter two properties are more intertwined with lobe-transitivity than the characterization of vertex-transitivity (for graphs of connectivity 1), scattered throughout are examples that illustrate some algebraic distinctions among these various properties.

2 Preliminaries

Throughout this article, the symbol \mathbb{N} denotes the set of positive integers. The symbols \mathbb{I}, \mathbb{J} , and \mathbb{K} , often subscripted, denote subsets of \mathbb{N} of the form $\{1, 2, \dots, n\}$ or the set \mathbb{N} itself; they appear as sets of indices. All graphs (and their valences) in this article are finite or countably infinite. The symbol $\delta_{i,j}$ (the so-called “Kronecker delta”) assumes the value 1 if $i = j$ and 0 if $i \neq j$. For a graph Λ and any subgroup $H \leq \text{Aut}(\Lambda)$, the set of orbits of H acting on $V\Lambda$ is denoted by $\mathcal{O}(H)$.

The set of lobes of a graph Γ is denoted by $\mathcal{L}(\Gamma)$. We let $\{\mathcal{L}_k : k \in \mathbb{K}\}$ denote the partition of $\mathcal{L}(\Gamma)$ into isomorphism classes of lobes. For given $k \in \mathbb{K}$ and a lobe $\Lambda \in \mathcal{L}_k$, we let $\mathcal{O}(\text{Aut}(\Lambda)) = \{(V\Lambda)_j : j \in \mathbb{J}_k\}$, and we understand that if $\sigma : \Lambda \rightarrow \Theta$ is an isomorphism between lobes in \mathcal{L}_k , then $\sigma((V\Lambda)_j) = (V\Theta)_j$ for all $j \in \mathbb{J}_k$. Finally, for each $k \in \mathbb{K}$ and $j \in \mathbb{J}_k$, we define the function $\tau_j^{(k)} : V\Gamma \rightarrow \mathbb{N}$ by

$$\tau_j^{(k)}(v) = |\{\Lambda \in \mathcal{L}_k : v \in (V\Lambda)_j\}|. \tag{2.1}$$

For $\Lambda_0 \in \mathcal{L}(\Gamma)$ and $n \in \mathbb{N}$, we recursively define the subgraphs

$$\begin{aligned} \Gamma_0(\Lambda_0) &= \Lambda_0, \\ \Gamma_{n+1}(\Lambda_0) &= \bigcup \{\Lambda \in \mathcal{L}(\Gamma) : V\Lambda \cap V\Gamma_n(\Lambda_0) \neq \emptyset\}. \end{aligned}$$

Lemma 2.1 ([5, Lemma 3.1]). *Let $\Lambda, \Theta \in \mathcal{L}(\Gamma)$ and let $n \in \mathbb{N}$. If for each $k \in \mathbb{K}$ and $j \in \mathbb{J}_k$, the function $\tau_j^{(k)}$ is constant on $V\Gamma$, then any isomorphism $\sigma_n : \Gamma_n(\Lambda) \rightarrow \Gamma_n(\Theta)$ admits an extension to an automorphism $\sigma \in \text{Aut}(\Gamma)$.*

²For a short algebraic proof that all 1-ended planar graphs with primitive automorphism group are biconnected, see [8].

This lemma was used in [5] to prove the following characterization of vertex-transitive graphs of connectivity 1.

Theorem 2.2 ([5, Theorem 3.2]). *Let Γ be a graph of connectivity 1. A necessary and sufficient condition for Γ to be vertex-transitive is that all the functions $\tau_j^{(k)}$ be constant on $V\Gamma$.*

Notation. When all the lobes of the graph Γ are pairwise isomorphic, that is, the index set \mathbb{K} has but one element, then in Equation (2.1) the index k is suppressed; we simply replace \mathbb{J}_k by \mathbb{J} and $\tau_j^{(k)}$ by τ_j .

3 Lobe-transitivity

Let Γ be a graph of connectivity 1. It is immediate from the above definitions that the edge-sets of the lobes of Γ are blocks of imprimitivity of the group $\text{Aut}(\Gamma)$ acting on $E\Gamma$. Hence any automorphism of Γ must map lobes onto lobes, and therefore, if $\text{Aut}(\Gamma)$ is to act transitively on $\mathcal{L}(\Gamma)$, then all the lobes of Γ must be pairwise isomorphic. However, pairwise-isomorphism of the lobes alone is not sufficient for lobe-transitivity, even when every vertex of Γ lies in the same number of lobes.

Let us first dispense with trees; the proof is elementary and hence omitted.

Proposition 3.1. *A finite or countable tree is lobe-transitive (and simultaneously, edge-transitive) if and only if there exist $n_1, n_2 \in \mathbb{N} \cup \{\aleph_0\}$ such that every edge has one incident vertex of valence n_1 and the other of valence n_2 . If $n_1 = n_2$, the tree is also arc-transitive.*

For graphs of connectivity 1 other than trees, we have the following characterization of lobe-transitivity.

Theorem 3.2. *Let Γ be a graph of connectivity 1, and let Λ_0 be an arbitrary lobe of Γ . Let $\{P_i : i \in \mathbb{I}\}$ be the set of orbits of $\text{Aut}(\Gamma)$, and let $\mathcal{Q} = \{Q_j : j \in \mathbb{J}\}$ be the set of those orbits of the stabilizer in $\text{Aut}(\Gamma)$ of Λ_0 that are contained in Λ_0 . Then necessary and sufficient conditions for the graph Γ to be lobe-transitive are:*

(1) *For each lobe $\Lambda \in \mathcal{L}(\Gamma)$, there exists an isomorphism $\sigma_\Lambda : \Lambda_0 \rightarrow \Lambda$.*

(2) *For each $j \in \mathbb{J}$, there exists a function $\tau_j : V\Gamma \rightarrow \mathbb{N} \cup \{0, \aleph_0\}$ such that*

(a) *for all $v \in V\Gamma$,*

$$\tau_j(v) = |\{\Lambda \in \mathcal{L}(\Gamma) : v \in \sigma_\Lambda(Q_j)\}| \tag{3.1}$$

and

(b) *for each $i \in \mathbb{I}$, τ_j is constant on P_i and is nonzero if and only if $Q_j \subset P_i$.*

Proof. (Necessity) Suppose that Γ is lobe-transitive. For each lobe $\Lambda \in \mathcal{L}(\Gamma)$, there is an automorphism $\bar{\sigma}_\Lambda \in \text{Aut}(\Gamma)$ that maps the fixed lobe Λ_0 onto Λ . The restriction to Λ_0 of $\bar{\sigma}_\Lambda$ is an isomorphism $\sigma_\Lambda : \Lambda_0 \rightarrow \Lambda$ that satisfies condition (1).

For any lobe Λ , an automorphism $\alpha \in \text{Aut}(\Gamma)$ is in the stabilizer of Λ if and only if $\bar{\sigma}_\Lambda^{-1}\alpha\bar{\sigma}_\Lambda$ is in the stabilizer of Λ_0 . It follows that the partition $\{\sigma_\Lambda(Q_j) : j \in \mathbb{J}\}$ of $V\Lambda$ is the set of orbits of the stabilizer of Λ that are contained in Λ . Furthermore, since the stabilizer of Λ_0 is a subgroup of $\text{Aut}(\Gamma)$, the partition $\{\sigma_\Lambda(Q_j) : j \in \mathbb{J}\}$ of $V\Lambda$

refines the partition $\{P_i \cap V\Lambda : i \in \mathbb{I}\}$. If for some indices i and j , the vertex v satisfies $v \in \sigma_\Lambda(Q_j) \subset P_i$, then for any lobe Θ , the vertex $\sigma_\Theta\sigma_\Lambda^{-1}(v)$ lies in $P_i \cap \sigma_\Theta(Q_j)$. This implies that, for all $j \in \mathbb{J}$, the function τ_j as given in Equation (3.1) is well-defined and constant on P_i .

Suppose that for an arbitrary index $i \in \mathbb{I}$, the vertex v lies in P_i . Since by Equation (3.1), $\tau_j(v)$ counts for each $j \in \mathbb{J}$ the number of lobes Λ such that v lies in $\sigma_\Lambda(Q_j)$, it follows that $\tau_j(v)$ is positive exactly when $\sigma_\Lambda^{-1}(v) \in Q_j \subset P_i$ holds, concluding the proof of condition (2.b).

(Sufficiency) Assume conditions (1) and (2). To prove that Γ is lobe-transitive, it suffices to prove that every isomorphism $\sigma_\Theta : \Lambda_0 \rightarrow \Theta$ is extendable to an automorphism of Γ . (Note that in this direction of the proof, σ_Θ is not presumed to be the restriction to Λ_0 of an automorphism $\bar{\sigma}_\Theta \in \text{Aut}(\Gamma)$ but, in fact, it is.)

Fix a lobe Θ_0 and a vertex $v \in V\Lambda_0$ and let $w = \sigma_{\Theta_0}(v)$. For some $j \in \mathbb{J}$, the vertex v lies in Q_j , and so $w \in \sigma_{\Theta_0}(Q_j)$. Since both $\tau_j(v)$ and $\tau_j(w)$ are therefore positive, both v and w lie in the same orbit P_i of $\text{Aut}(\Gamma)$ by condition (2.b). Furthermore, since τ_j is constant on P_i for each $j \in \mathbb{J}$, there exists a bijection β_j from the set of lobes Λ such that $v \in \sigma_\Lambda(Q_j)$ onto the set of lobes Θ such that $w \in \sigma_\Theta(Q_j)$. Let Λ_1 be a lobe in the former set, and let $\Theta_1 = \beta_j(\Lambda_1)$.

Although v and w lie in the images of the same orbit Q_j in lobes Λ_1 and Θ_1 , respectively, the vertices $\sigma_{\Lambda_1}^{-1}(v)$ and $\sigma_{\Theta_1}^{-1}(w)$ need not be the same vertex of Λ_0 . However, since both vertices lie in the same orbit Q_j of Λ_0 , there exists an automorphism $\alpha \in \text{Aut}(\Lambda_0)$ such that $\alpha\sigma_{\Lambda_1}^{-1}(v) = \sigma_{\Theta_1}^{-1}(w)$. Then $\sigma_{\Theta_1}\alpha\sigma_{\Lambda_1}^{-1}$ is an isomorphism from Λ_1 onto Θ_1 that maps v onto w and therefore agrees with σ_{Θ_0} at the vertex v common to Λ_0 and Λ_1 .

The amalgamation of $\sigma_{\Theta_1}\alpha\sigma_{\Lambda_1}^{-1}$ with σ_{Θ_0} is an isomorphism from $\Lambda_0 \cup \Lambda_1$ to $\Theta_0 \cup \Theta_1$. By repeating this same technique, we can extend σ_{Θ_0} to all lobes adjacent to Λ_0 and then to all of their adjacent lobes and inductively to all of Γ . □

Example 3.3. Suppose that the lobes of Γ are copies of some biconnected, vertex-transitive graph and that every vertex of Γ is incident with exactly m lobes where $m \geq 2$. By Theorem 2.2, Γ is vertex-transitive. By Theorem 3.2, Γ is lobe-transitive, with $\mathcal{O}(\text{Aut}(\Gamma))$ being the trivial partition (with just one big cell P_1). Also $|\mathbb{J}| = 1$ and $\tau_1(v) = m$ for all $v \in V\Gamma$.

Remark 3.4. There exists a “degenerate” family of lobe-transitive graphs Γ of connectivity 1 that have but a single cut-vertex. For some cardinal $\aleph \geq 2$, consider a collection of \aleph copies of a biconnected graph Λ_0 , and let $v_0 \in V\Lambda_0$. Let $\sigma_\Lambda : \Lambda_0 \rightarrow \Lambda$ be an isomorphism as in Theorem 3.2, and let $\sigma_\Lambda(v_0) = v_\Lambda$ for each copy Λ of Λ_0 in the collection. We obtain Γ by identifying v_0 and all the vertices v_Λ and naming the new amalgamated vertex w , which forms a singleton orbit $\{w\}$ of $\text{Aut}(\Gamma)$. Clearly Γ is lobe-transitive and w is its unique cut-vertex. If Λ_0 has finite diameter, then Γ has zero ends (see [3]) when Λ_0 is finite and has \aleph ends when Λ_0 is infinite; if Λ_0 has infinite diameter, then Γ has at least \aleph ends. Other than the graphs just described, all countable lobe-transitive graphs of connectivity 1 are “tree-like” with \aleph_0 cut vertices and either 2 or 2^{\aleph_0} ends.

Theorem 3.5. Let Λ_0 be any biconnected graph. Let $H \leq \text{Aut}(\Lambda_0)$, let $\mathcal{Q} = \mathcal{O}(H) = \{Q_j : j \in \mathbb{J}\}$, and let $\mathcal{R} = \{R_k : k \in \mathbb{K}\}$ be a partition of $V\Lambda_0$ refined by \mathcal{Q} . For each $k \in \mathbb{K}$, let the function $\mu_k : \mathbb{J} \rightarrow \mathbb{N} \cup \{0, \aleph_0\}$ satisfy $\mu_k(j) > 0$ if and only if $Q_j \subseteq R_k$ and (to avoid the triviality of a single lobe) $\sum_{j \in \mathbb{J}} \mu_k(j) \geq 2$ for at least one $k \in \mathbb{K}$. Then there exists (up to isomorphism) a unique lobe-transitive graph Γ of connectivity 1 such that

- (1) for each lobe $\Lambda \in \mathcal{L}(\Gamma)$, there exists an isomorphism $\sigma_\Lambda: \Lambda_0 \rightarrow \Lambda$;
- (2) for each vertex $v \in V\Gamma$ and each $j \in \mathbb{J}$, we have

$$\mu_k(j) = |\{\Lambda \in \mathcal{L}(\Gamma) : \sigma_\Lambda^{-1}(v) \in Q_j \subseteq R_k\}|.$$

Proof. Let Λ_0 , H , \mathcal{Q} , and μ_k be as postulated. Let $\Gamma_0 = \Lambda_0$ from which we construct Γ_1 as follows.

Let v be any vertex of Λ_0 . For some j, k , it must hold that $v \in Q_j \subseteq R_k$, and so $\mu_k(j) > 0$. For each ℓ such that $Q_\ell \subseteq R_k$, we postulate the existence of $\mu_k(\ell)$ copies Λ of Λ_0 (including Λ_0 itself when $\ell = j$) such that, if $\sigma_\Lambda: \Lambda_0 \rightarrow \Lambda$ is an isomorphism, then some vertex in $\sigma_\Lambda(Q_\ell)$ is identified with the vertex v . The graph Γ_1 is produced by repeating this process for each vertex of Λ_0 . We repeat this process starting at each vertex $w \in V\Gamma_1 \setminus V\Gamma_0$, the only notational change being that, if specifically $w \in \sigma_\Lambda(Q_{j'})$ for some $j' \in \mathbb{J}$, then we consider the subset $\sigma_\Lambda(Q_{j'})$ of $V\Lambda$ (instead of Q_j in Λ_0) to which w belongs. Thus we construct Γ_2 .

Inductively, suppose that Γ_n has been constructed for some $n \geq 2$. Let $w \in V\Gamma_n \setminus V\Gamma_{n-1}$, and so $w \in \sigma_\Lambda(Q_m)$ holds for some $m \in \mathbb{J}$ and a unique lobe $\Lambda \in \mathcal{L}(\Gamma_n) \setminus \mathcal{L}(\Gamma_{n-1})$. Supposing that $Q_m \subseteq R_k$, we postulate the existence of $\mu_k(m)$ new copies of Λ_0 that share only the vertex w with Γ_n according to the above identification. In this way we construct Γ_{n+1} . Finally, let $\Gamma = \bigcup_{n=0}^\infty \Gamma_n$.

It remains only to prove that Γ so-constructed is lobe-transitive. Let Θ be any lobe of Γ . By the above construction, all lobes of Γ are pairwise isomorphic, and so there exists an isomorphism $\sigma_\Theta: \Lambda_0 \rightarrow \Theta$. Starting with $\Gamma'_0 = \Theta$ and by using the technique in the proof of Sufficiency in Theorem 3.2, one constructs a sequence $\Gamma'_0, \Gamma'_1, \dots$ so that for all $n \in \mathbb{N}$, we have $\Gamma'_n \cong \Gamma_n$, and σ_Θ is extendable to an isomorphism from Γ_n to Γ'_n . Thus $\Gamma \cong \bigcup_{n=0}^\infty \Gamma'_n$, and σ_Θ can be extended to an automorphism of Γ . \square

Example 3.6. In the notation of Theorem 3.5, let Λ_0 be the 5-cycle with one chord as shown in Figure 1(a), and let $H = \text{Aut}(\Lambda)$, yielding the orbit partition $\{Q_1, Q_2, Q_3\}$ as indicated. Let $R_1 = Q_1 \cup Q_3$ and $R_2 = Q_2$, giving $\mathbb{J} = \{1, 2, 3\}$ and $\mathbb{K} = \{1, 2\}$. Define $\mu_1(1) = 3$, $\mu_1(3) = 1$, and $\mu_2(2) = 2$. Note that all other values of μ_1 and μ_2 must equal 0. Then Γ_1 is as seen in Figure 1(b).

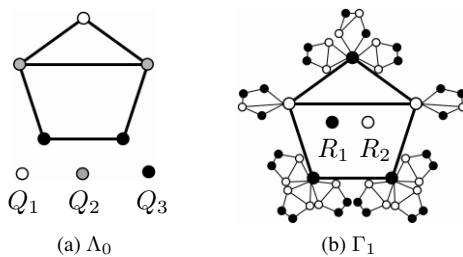


Figure 1: Λ_0 and Γ_1 from Example 3.6.

The pairs of conditions in Theorems 3.2 and 3.5 may appear alike, but there is a notable difference between them. This occurs when the arbitrarily chosen subgroup $H \leq \text{Aut}(\Lambda_0)$

of Theorem 3.5 is a *proper* subgroup of the stabilizer of Λ_0 in $\text{Aut}(\Gamma)$, where Γ is the graph constructed from Λ_0 and the functions μ_k of Theorem 3.5. We illustrate this distinction with following example.

Example 3.7. Our initial lobe Λ_0 is a copy of K_4 , with vertices labeled as in Figure 2(a), and so $\text{Aut}(\Lambda_0) \cong \text{Sym}(4)$ of order 24. We use Λ_0 to “build” the lobe-transitive graph Γ shown four times in Figure 2(b). The action on Λ_0 by the stabilizer of Λ_0 in $\text{Aut}(\Gamma)$ is the 4-element group $\langle g_1 \rangle \times \langle g_2 \rangle$ whose generators have cycle representation $g_1 = (v_1, v_2)$ and $g_2 = (v_3, v_4)$. The shadings of the vertices in the four depictions of Γ in Figure 2(b) correspond respectively to the four different subgroups of $\langle g_1 \rangle \times \langle g_2 \rangle$ described below. For the sake of simplicity, we assume $\mathcal{R} = \mathcal{Q}$.

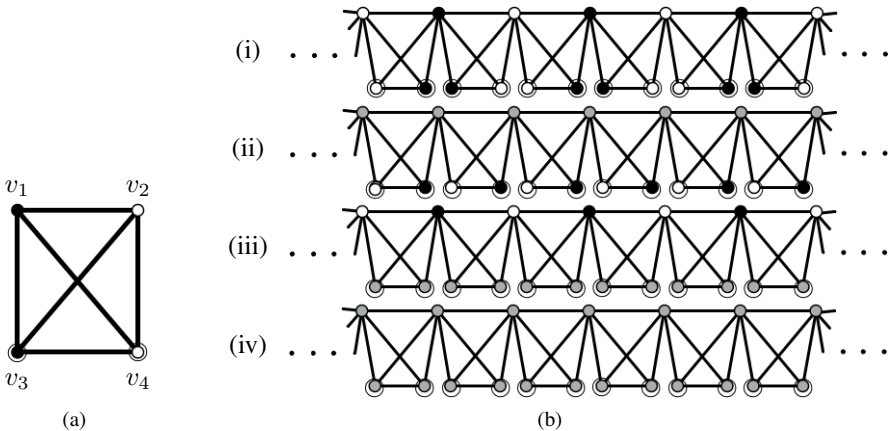


Figure 2: The clothesline graph.

- (i) H is the trivial group $\{\iota\}$. Thus H induces four orbits $Q_j = \{v_j\}$ for $j \in \{1, 2, 3, 4\}$. The functions μ_k are then given by $\mu_k(j) = 2\delta_{j,k}$ for $k \in \{1, 2\}$ and $\mu_k(j) = \delta_{j,k}$ for $k \in \{3, 4\}$.
- (ii) $H = \langle g_1 \rangle$. There are three orbits of H : $Q_1 = \{v_1, v_2\}$, $Q_2 = \{v_3\}$, and $Q_3 = \{v_4\}$, which give $\mu_1(1) = 2, \mu_2(2) = \mu_3(3) = 1$. All other functional values are zero.
- (iii) $H = \langle g_2 \rangle$. Again there are three orbits of H but not the same ones: $Q_1 = \{v_1\}$, $Q_2 = \{v_2\}$, and $Q_3 = \{v_3, v_4\}$. This gives $\mu_k(j) = 2\delta_{j,k}$ for $k \in \{1, 2\}$ and $\mu_3(j) = \delta_{j,3}$.
- (iv) $H = \langle g_1 \rangle \times \langle g_2 \rangle$. Now there are just two orbits: $Q_1 = \{v_1, v_2\}$ and $Q_2 = \{v_3, v_4\}$. Finally, $\mu_1(1) = 2, \mu_2(2) = 1$, and all other functional values are zero.

All four choices for H , the partition \mathcal{Q} , and the functions μ_k clearly yield the same lobe-transitive graph Γ of connectivity 1 by the construction of Theorem 3.5

4 Edge-transitivity

Lemma 4.1. *If Γ is an edge-transitive (respectively, arc-transitive) graph, then Γ is lobe-transitive and its lobes are also edge-transitive (respectively, arc-transitive).*

Proof. For $i = 1, 2$, let $\Theta_i \in \mathcal{L}(\Gamma)$, and let e_i be an edge (respectively, arc) of Θ_i . There exists an automorphism $\varphi \in \text{Aut}(\Gamma)$ such that $\varphi(e_1) = e_2$. Since φ maps cycles through e_1 onto cycles through e_2 , φ must map Θ_1 onto Θ_2 . If e_1 and e_2 lie in the same lobe Θ , then φ leaves Θ invariant, and so its restriction to Θ is an automorphism of Θ . \square

Theorem 4.2. *Let Γ be a graph of connectivity 1 with more than one lobe, and let $\Lambda \in \mathcal{L}(\Gamma)$. Necessary and sufficient conditions for Γ to be edge-transitive are the following:*

- (1) *The lobes of Γ are edge-transitive.*
- (2) *For each lobe $\Theta \in \mathcal{L}(\Gamma)$, there exists an isomorphism $\sigma_\Theta: \Lambda \rightarrow \Theta$.*
- (3) *Exactly one of the following descriptions of Γ holds:*
 - (a) *Both Γ and Λ are vertex-transitive, in which case every vertex is incident with the same number ≥ 2 of lobes.*
 - (b) *The graph Γ is vertex-transitive but Λ is not vertex-transitive, in which case Λ is bipartite with bipartition $\{Q_1, Q_2\}$, and there exist constants $m_1, m_2 \in \mathbb{N} \cup \{\aleph_0\}$ such that for $j = 1, 2$ and all $v \in V\Gamma$, it holds that*

$$m_j = |\{\Theta \in \mathcal{L}(\Gamma) : v \in \sigma_\Theta(Q_j)\}|.$$

- (c) *The graph Γ is not vertex-transitive, in which case Γ is bipartite with bipartition $\{P_1, P_2\}$ and there exist constants $m_1, m_2 \in \mathbb{N} \cup \{\aleph_0\}$, at least one of which is at least 2, such that for $i = 1, 2$, if $v \in P_i$, then*

$$|\{\Theta \in \mathcal{L}(\Gamma) : v \in \sigma_\Theta(P_j \cap V\Lambda)\}| = m_j \delta_{i,j}.$$

Proof. (Sufficiency) Assume all the conditions in the hypothesis and let $e_1, e_2 \in E\Gamma$ be arbitrary edges in lobes Λ_1 and Λ_2 , respectively. By condition (2), there exists an isomorphism $\sigma: \Lambda_1 \rightarrow \Lambda_2$. By condition (1), there exists an automorphism $\alpha \in \text{Aut}(\Lambda_2)$ such that $e_2 = \alpha\sigma(e_1)$. Each of the three cases of condition (3) is seen to satisfy the hypothesis of Lemma 2.1, implying that $\alpha\sigma: \Lambda_1 \rightarrow \Lambda_2$ is extendable to an isomorphism of Γ mapping e_1 to e_2 .

(Necessity) Suppose that Γ is an edge-transitive graph of connectivity 1. By Lemma 4.1, Γ is also lobe-transitive and its lobes are edge-transitive, proving condition (1). Condition (2), which establishes notation for the remainder of this proof, also follows from Lemma 4.1.

To prove (3), we continue the notation of Theorem 3.2 with \mathbb{I} being the index set for the set of orbits of $\text{Aut}(\Gamma)$ and \mathbb{J} being the index set for the orbits of the stabilizer of Λ that are contained in Λ . Since both Γ and all of its lobes are edge-transitive, $|\mathbb{I}|$ and $|\mathbb{J}|$ equal either 1 or 2.

If both Γ and Λ are vertex-transitive, then $|\mathbb{I}| = |\mathbb{J}| = 1$, and for every vertex $v \in V\Gamma$, $\tau_1(v) = m$ holds for some $m \geq 2$. This is case (3.a).

Since any odd cycle in Γ would be contained in a lobe of Γ , it holds that Γ is bipartite if and only if every lobe is bipartite. If either Γ or Λ is not vertex-transitive, then each – and hence both – are bipartite, and the sides of the bipartitions (whether or not they are entire orbits of the appropriate automorphism group) are blocks of imprimitivity systems. Let $\{P_1, P_2\}$ be the bipartition of $V\Gamma$, and so $\{P_1 \cap V\Theta, P_2 \cap V\Theta\}$ is the bipartition of

any lobe Θ . Equivalently, letting $\{Q_1, Q_2\}$ denote the bipartition of Λ , we have $P_i = \bigcup_{\Theta \in \mathcal{L}(\Gamma)} \sigma_\Theta(Q_i)$ for $i = 1, 2$.

Suppose that Γ is vertex-transitive but Λ is not, and so $|\mathbb{I}| = 1$ and $|\mathbb{J}| = 2$. By Theorem 2.2, there exist constants $m_1, m_2 \in \mathbb{N} \cup \{\aleph_0\}$ such that for all $v \in V\Gamma$ and $j \in \mathbb{J}$, we have $m_j = \tau_j(v) = |\{\Theta \in \mathcal{L}(\Gamma) : v \in \sigma_\Theta(Q_j)\}|$.

Finally, suppose that Γ is not vertex-transitive, and so P_1 and P_2 are the orbits of $\text{Aut}(\Gamma)$. Also, Λ is bipartite with bipartition $\{Q_1, Q_2\}$, where $Q_i = P_i \cap V\Lambda$. As no automorphism of Γ swaps P_1 with P_2 , no automorphism of Γ swaps Q_1 with Q_2 (even though Λ may be vertex-transitive!). Hence $|\mathbb{I}| = |\mathbb{J}| = 2$. Since Γ is lobe-transitive, it follows now from the “necessity” argument of Theorem 3.2 that, for $j = 1, 2$, the function τ_j satisfies the condition $\tau_j(v) > 0$ if and only if $v \in P_j$. That means that there exist constants $m_1, m_2 \in \mathbb{N} \cup \{\aleph_0\}$, at least one of which is greater than 1, such that, if $v \in P_i$, then $\tau_j(v) = m_j \delta_{i,j}$. \square

Example 4.3. Suppose in the notation of Theorem 4.2 that Γ is edge-transitive and Λ is the complete bipartite graph $K_{s,t}$ with $|Q_1| = s$ and $|Q_2| = t$. Suppose that every vertex of Γ is incident with exactly two lobes isomorphic to Λ . If $s = t$, then both Γ and Λ are vertex-transitive, and we have case (3.a) of the theorem. If $s \neq t$ and every vertex lies in one image of Q_1 and one image of Q_2 , then we have the situation of case (3.b). If again $s \neq t$ but each vertex lies in either two images of Q_1 or two images of Q_2 , then we have the situation described in case (3.c).

Remark 4.4. With regard to Example 4.3, we note that having $s = t$ does not assure vertex-transitivity of edge-transitive bipartite graphs. There exist edge-transitive, non-vertex-transitive, finite bipartite graphs where the two sides of the bipartition have the same size. Such graphs are called *semisymmetric*. The smallest such graph, on 20 vertices with valence 4, was found by J. Folkman [2], who also found several infinite families of semisymmetric graphs. Many more such families as well as forbidden values for s were determined by A. V. Ivanov [4].

Example 4.5. This simple example illustrates how the converse of Lemma 4.1 is false, even though the lobes themselves may be highly symmetric. Let Γ be a graph of connectivity 1 whose lobes are copies of the Petersen graph (which is 3-arc-transitive!). For each lobe Λ , let $(V\Lambda)_1$ and $(V\Lambda)_2$ denote the vertex sets of disjoint 5-cycles indexed “consistently,” i.e., if Λ and Θ share a vertex v , then $v \in (V\Lambda)_i \cap (V\Theta)_i$ for $i = 1$ or $i = 2$. (Observe that Γ is not bipartite.) For $i = 1, 2$ define $P_i = \bigcup_{\Theta \in \mathcal{L}(\Gamma)} (V\Theta)_i$, and suppose that each vertex in P_i belongs to exactly m_i lobes. The graph Γ is lobe-transitive by Theorem 3.2, and Γ is both vertex- and edge-transitive when $m_1 = m_2$, but Γ is neither vertex- nor edge-transitive when $m_1 \neq m_2$.

5 Arc-transitivity

Theorem 5.1. *Let Γ be a graph of connectivity 1. Necessary and sufficient conditions for Γ to be arc-transitive are the following:*

- (1) *The lobes of Γ are arc-transitive.*
- (2) *The lobes of Γ are pairwise isomorphic.*
- (3) *All vertices of Γ are incident with the same number of lobes.*

Proof. (Necessity) Suppose that Γ is arc-transitive. Conditions (1) and (2) follow from Lemma 4.1. Since arc-transitivity implies vertex-transitivity, condition (3) holds.

(Sufficiency) Assume that the three conditions hold. For $k = 1, 2$, let a_k be an arc of Γ , and let Θ_k be the lobe containing a_k . By condition (2), there exists an isomorphism $\sigma: \Theta_1 \rightarrow \Theta_2$. By condition (1), there exists an automorphism $\alpha \in \text{Aut}(\Theta_2)$ such that $\alpha\sigma(a_1) = a_2$. By condition (3), the functions τ_j of Equation (2.1) are constant on $V\Gamma$. (In fact, since the lobes are vertex-transitive, there is only one such function.) It now follows from Lemma 2.1 that $\alpha\sigma$ is extendable to all of Γ . \square

Remark 5.2. If conditions (1) and (3) of Theorem 5.1 were replaced by *the lobes are edge-transitive and Γ is vertex-transitive*, the *sufficiency* argument would fail. There exist finite graphs [1] and countably infinite graphs with polynomial growth rate [7] that are vertex- and edge-transitive but not arc-transitive. Let Λ denote such a graph, and consider a graph Γ whose lobes are isomorphic to Λ with the same number of lobes incident with every vertex. Then Γ itself is vertex- and edge-transitive but not arc-transitive.

The following proposition is elementary.

Proposition 5.3. *For all $k \geq 2$, the only k -arc-transitive graphs of connectivity 1 are trees of constant valence.*

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